

# Packrat Parsers Can Support Left Recursion

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# Packrat Parsing

- Bryan Ford's ICFP'02 "Functional Pearl"
- Memoization of intermediate results → linear parse times
  - Backtracking
  - Unlimited look-ahead
- No ambiguities...

# Ordered Choice

- The expression  $e_1 / e_2$  means:
  - try  $e_1$
  - if successful, return result
  - otherwise, **backtrack** and try  $e_2$
- Makes parser's behavior easy to understand

# Packrat Infestation!

- Dozens of implementations
  - Pappy, Rats!, LPeg, ...
- Used in lots of projects
  - Fortress, Matchete, ...
- We use them for program transformation (e.g., OMeta, CAT)



# Left Recursion

- Natural way to express syntax of left-associative operators
- Left recursive rules → left-associative ASTs

$\text{expr} ::= \text{expr} \text{ “-” } \text{number}$ $\quad \quad \quad / \quad \text{number}$
--

- **Problem:** top-down parsers do not support left recursion...

# ... but packrat parsers are different!

- Intermediate results stored in the parser's **memo table**
- Our paper:



A way to leverage the memo table to support left recursion

# Technical Contributions

- Algorithm for supporting left recursion
- Experimental results:
  - typical uses of left recursion supported in linear time
  - very little overhead for non-left-recursive rules
  - can parse heavily left-recursive subset of the Java grammar (as defined in the JLS)

# An Alternative Approach

- Rewrite left-recursive grammars
- Technique developed for CFGs, not fully understood in the context of ordered choice
- Pappy [Ford'02] and Rats! [Grimm'06] rewrite **directly left-recursive rules**
- **Indirectly left-recursive** grammars must be rewritten manually



# Outline

- Memoization in packrat parsers
- Leveraging memoization to support left recursion
- Further details
- Performance
- Related work

# Memoization

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{number} \text{ “-” } \text{expr} \\ \quad \quad \quad / \text{ number} \\ \text{stmt} ::= \text{expr} \text{ “;”} \\ \quad \quad \quad / \text{expr} \text{ “.”} \end{array} \right\}$

	0	1	2	3	4
expr					
number					
stmt					

# Memoization

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{number} \text{ “-” } \text{expr} \\ \quad \quad \quad / \text{ number} \\ \text{stmt} ::= \text{expr} \text{ “;”} \\ \quad \quad \quad / \text{expr} \text{ “.”} \end{array} \right\}$

	0	1	2	3	4
expr					
number					
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{number} \text{ “-” } \text{expr} \\ \quad \quad \quad / \text{ number} \\ \text{stmt} ::= \text{expr} \text{ “;”} \\ \quad \quad \quad / \text{expr} \text{ “.”} \end{array} \right\}$

	0	1	2	3	4
expr	?				
number					
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{number} \text{ “-” } \text{expr} \\ \quad \quad \quad / \text{ number} \\ \text{stmt} ::= \text{expr} \text{ “;”} \\ \quad \quad \quad / \text{expr} \text{ “.”} \end{array} \right\}$

	0	1	2	3	4
expr	?				
number					
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{number} \text{ “-” } \text{expr} \\ \quad \quad \quad / \text{ number} \\ \text{stmt} ::= \text{expr} \text{ “;”} \\ \quad \quad \quad / \text{expr} \text{ “.”} \end{array} \right\}$

	0	1	2	3	4
expr	?				
number					
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{number} \text{ “-” } \text{expr} \\ \quad \quad \quad / \text{ number} \\ \text{stmt} ::= \text{expr} \text{ “;”} \\ \quad \quad \quad / \text{expr} \text{ “.”} \end{array} \right\}$

	0	1	2	3	4
expr	?		?		
number	1				
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar = {  
 $\text{expr} ::= \text{number} \text{“-”} \text{expr}$   
 $\quad \quad \quad / \text{number}$   
 $\text{stmt} ::= \text{expr} \text{“;”}$   
 $\quad \quad \quad / \text{expr} \text{“.”}$

	0	1	2	3	4
expr	?		?		
number	1		2		
stmt	?				



# Memoization

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{number} \text{ “-” } \text{expr} \\ \quad \quad \quad / \text{ number} \\ \text{stmt} ::= \text{expr} \text{ “;”} \\ \quad \quad \quad / \text{expr} \text{ “.”} \end{array} \right\}$

	0	1	2	3	4
expr	?		?		
number	1		2		
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{number} \text{ “-” } \text{expr} \\ \quad \quad \quad / \text{ number} \\ \text{stmt} ::= \text{expr} \text{ “;”} \\ \quad \quad \quad / \text{expr} \text{ “.”} \end{array} \right\}$

	0	1	2	3	4
expr	?		?		?
number	1		2		
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar = {  
 $\text{expr} ::= \text{number} \text{“-”} \text{expr}$   
 $\quad \quad \quad / \text{number}$   
 $\text{stmt} ::= \text{expr} \text{“;”}$   
 $\quad \quad \quad / \text{expr} \text{“.”}$

	0	1	2	3	4
expr	?		?		?
number	1		2		3
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{number} \text{ “-” } \text{expr} \\ \quad \quad \quad / \text{ number} \\ \text{stmt} ::= \text{expr} \text{ “;”} \\ \quad \quad \quad / \text{expr} \text{ “.”} \end{array} \right\}$

	0	1	2	3	4
expr	?		?		?
number	1		2		3
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar = {  
 $\text{expr} ::= \text{number} \text{“-”} \text{expr}$   
 $\quad \quad \quad / \text{number}$   
 $\text{stmt} ::= \text{expr} \text{“;”}$   
 $\quad \quad \quad / \text{expr} \text{“.”}$

	0	1	2	3	4
expr	?		?		?
number	1		2		3
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{number} \text{ “-” } \text{expr} \\ \quad \quad \quad / \text{ number} \\ \text{stmt} ::= \text{expr} \text{ “;”} \\ \quad \quad \quad / \text{expr} \text{ “.”} \end{array} \right\}$

	0	1	2	3	4
expr	?		?		3
number	1		2		3
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar = {  
 $\text{expr} ::= \text{number} \text{“-”} \text{expr}$   
 $\quad \quad \quad / \text{number}$   
 $\text{stmt} ::= \text{expr} \text{“;”}$   
 $\quad \quad \quad / \text{expr} \text{“.”}$

	0	1	2	3	4
expr	?		(- 2 3)		3
number	1		2		3
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar = {  
 $\text{expr} ::= \text{number} \text{“-”} \text{expr}$   
 $\quad \quad \quad / \text{number}$   
 $\text{stmt} ::= \text{expr} \text{“;”}$   
 $\quad \quad \quad / \text{expr} \text{“.”}$

	0	1	2	3	4
expr	(- 1 (- 2 3))		(- 2 3)		3
number	1		2		3
stmt	?				



# Memoization

Input  
1-2-3.  
^

Grammar = {  
 $\text{expr} ::= \text{number} \text{“-”} \text{expr}$   
 $\quad \quad \quad / \text{number}$   
 $\text{stmt} ::= \text{expr} \text{“;”}$   
 $\quad \quad \quad / \text{expr} \text{“.”}$

	0	1	2	3	4
expr	(- 1 (- 2 3))		(- 2 3)		3
number	1		2		3
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar = {  
 $\text{expr} ::= \text{number} \text{“-”} \text{expr}$   
 $\quad \quad \quad / \text{number}$   
 $\text{stmt} ::= \text{expr} \text{“;”}$   
 $\quad \quad \quad / \text{expr} \text{“.”}$

	0	1	2	3	4
expr	(- 1 (- 2 3))		(- 2 3)		3
number	1		2		3
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{number} \text{ “-” } \text{expr} \\ \quad \quad \quad / \text{ number} \\ \text{stmt} ::= \text{expr} \text{ “;”} \\ \quad \quad \quad / \text{expr} \text{ “.”} \end{array} \right\}$

	0	1	2	3	4
expr	(- 1 (- 2 3))		(- 2 3)		3
number	1		2		3
stmt	?				

# Memoization

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{number} \text{ “-” } \text{expr} \\ \quad \quad \quad / \text{ number} \\ \text{stmt} ::= \text{expr} \text{ “;”} \\ \quad \quad \quad / \text{expr} \text{ “.”} \end{array} \right\}$

	0	1	2	3	4
expr	(- 1 (- 2 3))		(- 2 3)		3
number	1		2		3
stmt	(- 1 (- 2 3))				

$\text{expr} ::= \text{number} \text{“-”} \text{expr}$   
 $\quad \quad \quad / \text{number}$

$(- 1 (- 2 3))$

expr ::= number “-” expr  
/ number

(- 1 (- 2 3))

expr ::= expr “-” number  
/ number

(- (- 1 2) 3)

# Left-recursion = trouble

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{expr} \text{“-”} \text{number} \\ \quad \quad \quad / \text{number} \\ \text{stmt} ::= \text{expr} \text{“;”} \\ \quad \quad \quad / \text{expr} \text{“.”} \end{array} \right\}$

	0	1	2	3	4
expr					
number					
stmt					

# Left-recursion = trouble

Input  
1-2-3.  
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Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{expr} \text{“-”} \text{number} \\ \quad \quad \quad / \text{number} \\ \text{stmt} ::= \text{expr} \text{“;”} \\ \quad \quad \quad / \text{expr} \text{“.”} \end{array} \right\}$

	0	1	2	3	4
expr					
number					
stmt	?				



# Left-recursion = trouble

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	0	1	2	3	4
expr	?				
number					
stmt	?				

# Left-recursion = trouble

Input  
1-2-3.  
^

Grammar =  $\left\{ \begin{array}{l} \text{expr} ::= \text{expr} \text{“-”} \text{number} \\ \quad \quad \quad / \text{number} \\ \text{stmt} ::= \text{expr} \text{“;”} \\ \quad \quad \quad / \text{expr} \text{“.”} \end{array} \right\}$

	0	1	2	3	4
expr	?				
number					
stmt	?				

**Infinite loop!**



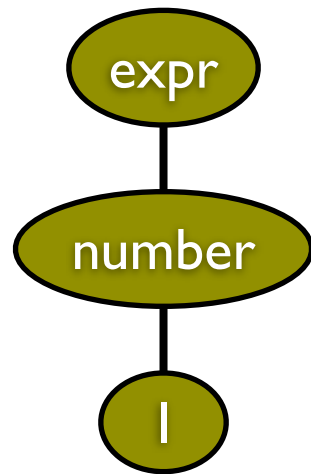
**Warning:  
Super-Duper  
Important  
Slide**

1-2-3

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$   
 $\quad \quad \quad / \text{number}$

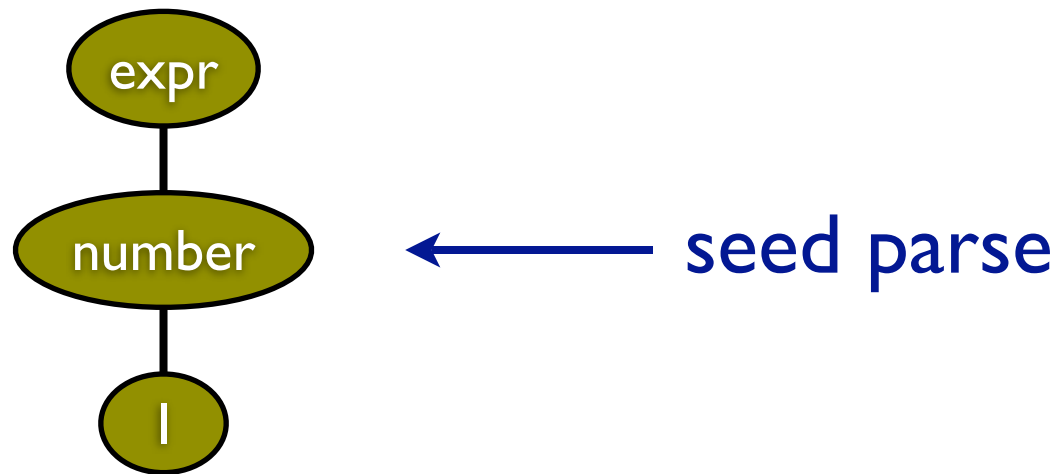
1-2-3

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$   
 $\quad \quad \quad / \text{number}$



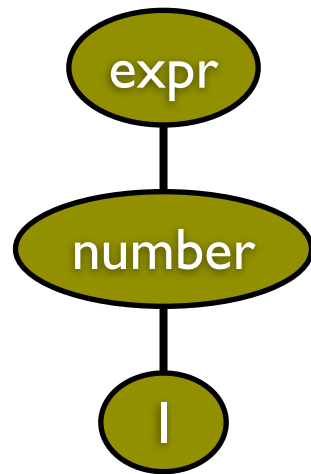
1-2-3

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$   
 $\quad \quad \quad / \text{number}$



1-2-3

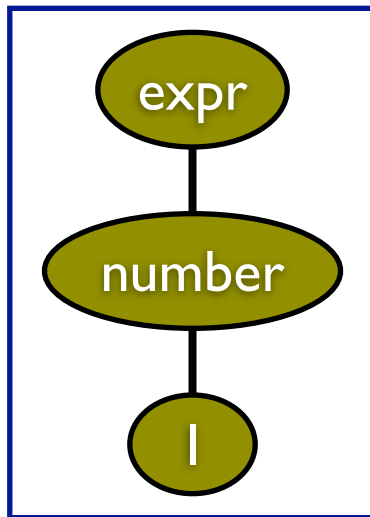
$\text{expr} ::= \text{expr} \text{“-”} \text{number}$   
 $\quad \quad \quad / \text{number}$





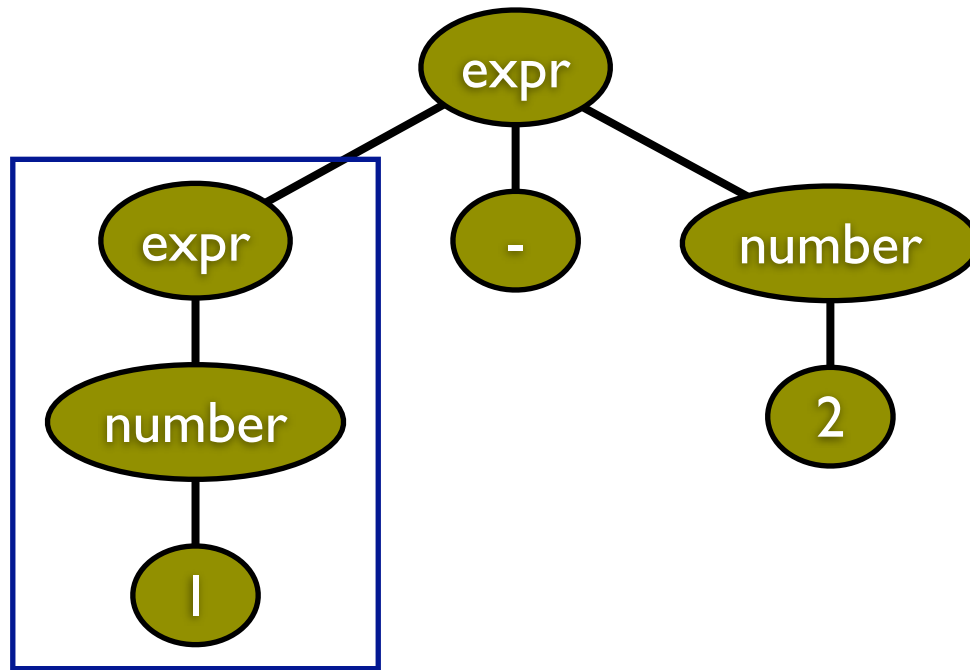
1-2-3

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$   
 $\quad \quad \quad / \text{number}$



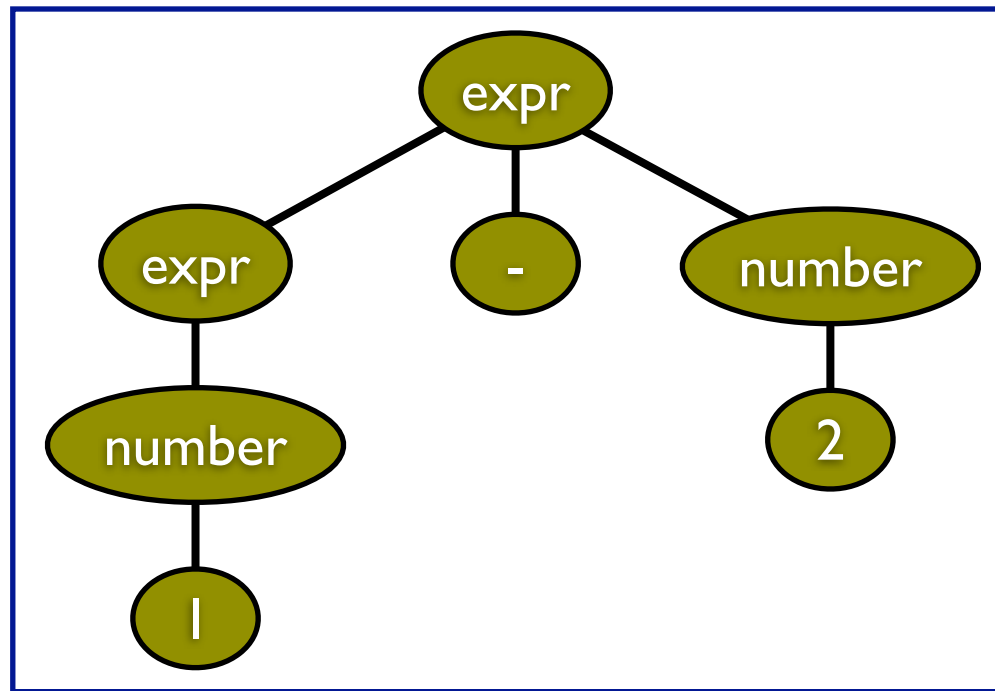
1-2-3

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$   
 $\quad \quad \quad / \text{number}$



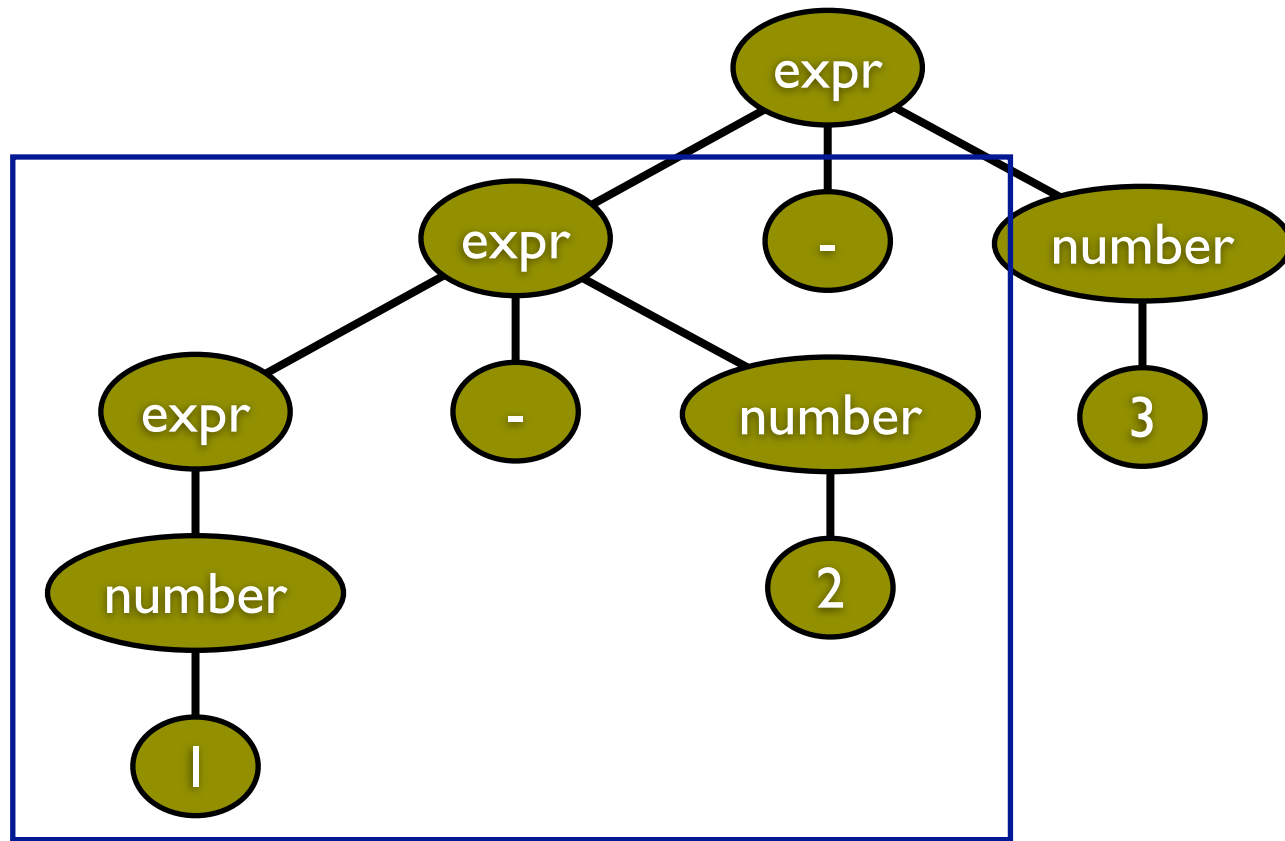
1-2-3

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$   
 $\quad \quad \quad / \text{number}$



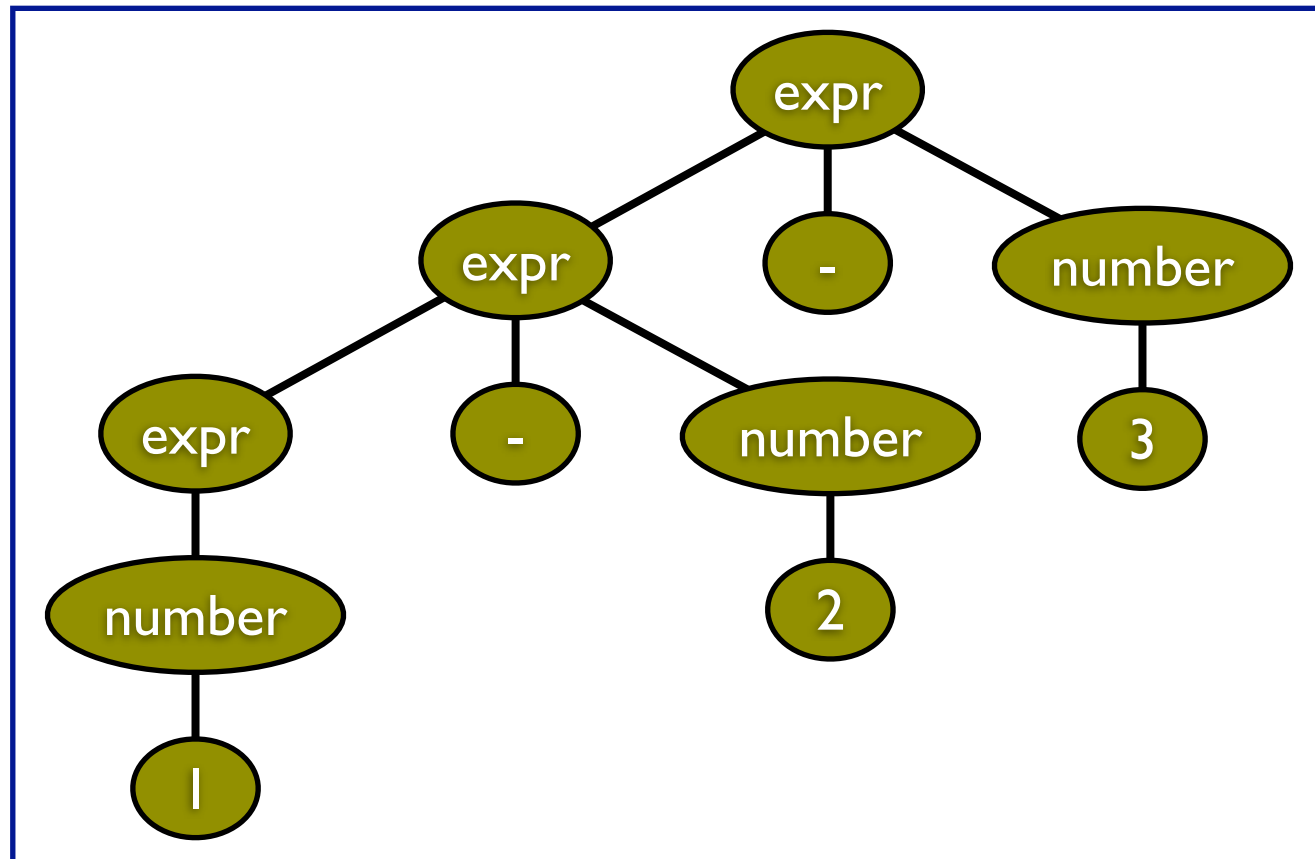
1-2-3

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$   
 $\quad \quad \quad / \text{number}$



1-2-3

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$   
 $\quad \quad \quad / \text{number}$



# Finding the Seed

Input = 1-2-3  
          <sup>^</sup>

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$ $\quad \quad \quad / \text{number}$
--

	0	1	2	3	4	5
expr						
number						

# Finding the Seed

Input = 1-2-3  
          <sup>^</sup>

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$ $\quad \quad \quad / \text{number}$
--

	0	1	2	3	4	5
expr	<b>FAIL</b>					
number						

# Finding the Seed

Input = 1-2-3  
          <sup>^</sup>

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$ $\quad \quad \quad / \text{number}$
--

	0	1	2	3	4	5
expr	<b>FAIL</b>					
number						



# Finding the Seed

Input = 1-2-3  
          <sup>^</sup>

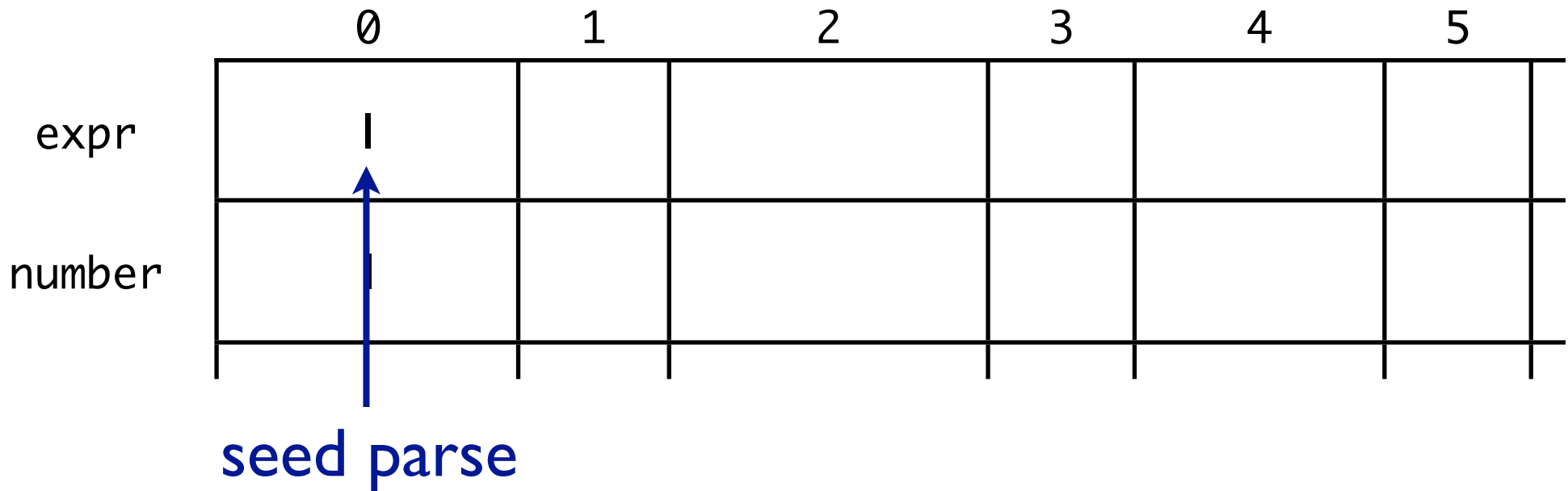
expr ::= expr “-” number  
      / number

	0	1	2	3	4	5
expr						
number						

# Finding the Seed

Input = 1-2-3  
          ^

$\text{expr} ::= \text{expr} \text{"-"} \text{number}$ $\quad \quad \quad / \text{number}$
---



# Finding the Seed

Input = 1-2-3  
          <sup>^</sup>

$\text{expr} ::= \text{expr} \text{"-"} \text{number}$ $\quad \quad \quad / \text{number}$
---

	0	1	2	3	4	5
expr						
number						

# Growing the Seed

Input = 1-2-3  
          <sup>^</sup>

$\text{expr} ::= \text{expr} \text{"-"} \text{number}$ $\quad \quad \quad / \text{number}$
---

	0	1	2	3	4	5
expr						
number						

# Growing the Seed

Input = 1-2-3  
          <sup>^</sup>

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$ $\quad \quad \quad / \text{number}$
--

	0	1	2	3	4	5
expr						
number						

# Growing the Seed

Input = 1-2-3  
          <sup>^</sup>

$\text{expr} ::= \text{expr} \text{"-"} \text{number}$ $\quad \quad \quad / \text{number}$
---

	0	1	2	3	4	5
expr						
number						

# Growing the Seed

Input = 1-2-3  
          <sup>^</sup>

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$ $\quad \quad \quad / \text{number}$
---

	0	1	2	3	4	5
expr						
number						

# Growing the Seed

Input = 1-2-3  
          <sup>^</sup>

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$ $\quad \quad \quad / \text{number}$
---

	0	1	2	3	4	5
expr	1					
number	1		2			



# Growing the Seed

Input = 1-2-3  
          <sup>^</sup>

expr ::= expr “-” number  
      / number

	0	1	2	3	4	5
expr	(-   2)					
number	1		2			

# Growing the Seed

Input = 1-2-3  
          <sup>^</sup>

$\text{expr} ::= \text{expr} \text{“-”} \text{number}$ $\quad \quad \quad / \text{number}$
---

	0	1	2	3	4	5
expr	(-   2)					
number	1		2			

# Growing the Seed

Input = 1-2-3  
          <sup>^</sup>

expr ::= expr “-” number  
      / number

	0	1	2	3	4	5
expr	(-   2)					
number	1		2			

# Growing the Seed

Input = 1-2-3  
          <sup>^</sup>

expr ::= expr “-” number  
      / number

	0	1	2	3	4	5
expr	(-   2)					
number	1		2			

# Growing the Seed

Input = 1-2-3<sub>^</sub>

expr ::= expr “-” number  
/ number

	0	1	2	3	4	5
expr	(-   2)					
number	1		2		3	

# Growing the Seed

Input = 1-2-3  
          <sup>^</sup>

expr ::= expr “-” number  
      / number

	0	1	2	3	4	5
expr	(- (- 1 2) 3)					
number	1		2		3	

# Growing the Seed

Input = 1-2-3<sup>^</sup>

expr ::= expr “-” number  
/ number

	0	1	2	3	4	5
expr	(- (- 1 2) 3)					
number	1		2		3	

# Growing the Seed

Input = 1-2-3<sub>^</sub>

expr ::= expr “-” number  
/ number

	0	1	2	3	4	5
expr	(- (- 1 2) 3)					
number	1		2		3	



# Other Aspects of the Algorithm

- Avoiding unnecessary work for non-left-recursive rules
- Supporting indirect left recursion
- See paper for details

# Performance (I)

- Experimental results:
  - Our approach supports typical uses of left recursion in linear time
  - It introduces very little overhead for non-left-recursive rules
  - Left recursion faster than right recursion (w/o tail call optimization)

# Performance (2)

- **Bad news:** possibly super-linear parse times
- **Good news:** only for contrived grammars

ones	::=	ones	“1”
		/	“1”
start	::=	ones	“2”
		/	“1” start

1111111111  
—————  
—————  
—————  
—————  
—————  
—————  
—————  
—————

# Related Work (I)

- [Frost & Hafiz'06]
  - can support left recursion by limiting otherwise infinite left recursion to  $N-1$  levels
  - works for any top-down parser, but
    - must know length of input stream
    - $O(n^4)$

# Related Work (2)

- [Johnson'95]: technique for building parsers for CFGs
  - based on memoization and CPS
  - left recursion support, polynomial parse times

# Related Work (3)

- Katahdin [Seaton'07]
  - language w/ extensible syntax
  - supports rules **annotated as left-recursive** using similar iterative process
  - does not support indirect left recursion

# Conclusion

- Packrat parsers can support left recursion
  - w/o left recursion elimination
  - usually in linear time
  - straightforward implementation (see paper)

**The End**



**“Just-in-case”**

**Slides**

# Ford's Transformation (I)

$\text{number} ::= n:\text{number } d:\text{digit} \rightarrow n * 10 + d$
$\quad \quad \quad / d:\text{digit} \quad \quad \quad \rightarrow d$



$\text{number} ::= d:\text{digit } f:\text{numberTail} \rightarrow f(d)$
$\text{numberTail} ::= d:\text{digit } \text{numberTail}:f \rightarrow \lambda n.f(n * 10 + d)$
$\quad \quad \quad / \text{empty} \quad \quad \quad \rightarrow \lambda x.x$

# Ford's Transformation (2)

Input = 123

	0	1	2	3
digit	1, pos'=1	2, pos'=2	3, pos'=3	<b>FAIL</b>
number	123, pos'=3	23, pos'=3	3, pos'=3	<b>FAIL</b>
numberTail			$\lambda n.(\lambda x.x)(n * 10 + 3),$ pos'=3	$\lambda x.x,$ pos'=3



$\lambda n.(\lambda n.(\lambda x.x)(n * 10 + 3))(n * 10 + 2),$   
pos'=3

# Ford's Transformation (3)

- From Bryan Ford's thesis:
  - *“As long as the computation of each cell looks up only a limited number of previously-recorded cells in the matrix and completes in constant time, the parsing process as a whole completes in linear time.”*
- At  $\text{pos}=i$ , the function returned by `numberTail` could perform  $n-i$  additions and multiplications
- So computation of `number` takes  $O(n)$ 
  - **Violates constant time stipulation!!**